
From the Selected Works of Lester G Telser

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Making Nöther Matrics

Lester G Telser

Circuits Using Nöther Matrices [NM]

L. G.Telser

An arrow, $a[i,j]$, in an m -polygon goes from vertex i , its source, to vertex j , its destination, $i \leftrightarrow j$. It is not two dimensional in a Nöther Algebra. It is a non-zero term in an $s \times s$ matrix whose terms are zero with this one exception, the term in row i , column j equals 1. In my work arrows always belong to circuits. A circuit with s arrows has a given relative order such that the destination of arrow $a[i,j]$, is the source of the next arrow, $a[j,k]$. The source of the arrow $a[j,k]$ is the destination of its predecessor, $a[i,j]$. Arrows in a circuit describe a round trip. Every arrow in an m -polygon belongs to at least one circuit. They never are alone. To show this the symbol for an arrow uses a special notation. Arrow $a[i,j]$ in an s -circuit has the special symbol $G[i,j,s]$.

A matrix in N -algebra has s rows and columns to represent a one-way arrow in an s -circuit. Matrix distinguishes it from matrix. An s -circuit is the sequence

$$\{a[i_1, j_1], a[i_1, j_2], a[i_2, j_3], \dots, a[i_s, j_1]\}.$$

As a matrix it becomes the product

$G[i_1, j_1, s] \cdot G[i_1, j_2, s] \cdot G[i_2, j_3, s] \cdot \dots \cdot G[i_s, j_s, s] = G[i_1, i_1, s]$, an $s \times s$ matrix with one term =1 at row i_1 , column i_1 , a diagonal. All its other terms are 0. Because the same s appears in each $G[.]$, in most cases s is omitted since it is clear from the context. This algebra uses an $s \times s$ matrix to identify an arrow in the s -circuit.

1. Simple Circuits

In[]*:= **Array**[0 &, {3, 3}]

Out[]*= { {0, 0, 0}, {0, 0, 0}, {0, 0, 0} }

```
In[•]:= MatrixForm[Array[0 &, {3, 3}]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

```
In[•]:= G[1, 2] := {{0, 1, 0}, {0, 0, 0}, {0, 0, 0}}
```

```
In[•]:= MatrixForm[G[1, 2]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

```
In[•]:= G[2, 3] := {{0, 0, 0}, {0, 0, 1}, {0, 0, 0}}
```

```
In[•]:= MatrixForm[G[2, 3]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$$

```
In[•]:= G[3, 1] := {{0, 0, 0}, {0, 0, 0}, {1, 0, 0}}
```

```
In[•]:= MatrixForm[G[3, 1]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 1 & 0 & 0 \end{pmatrix}$$

The symbol for the simple circuit (1, 2, 3) is the product of 3 matrices, one matrix for each one-way arrows in the simple 3-circuit. The term indicates a round trip starting from vertex 1. It appears as the first term on the diagonal of the N-matrix that shows the product of the 3 one-way arrows in their matrices.

```
In[•]:= MatrixForm[G[1, 2].G[2, 3].G[3, 1]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

```
In[•]:= G[1, 1] := {{1, 0, 0}, {0, 0, 0}, {0, 0, 0}}
```

```
In[•]:= MatrixForm[G[1, 1]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

```
In[•]:= G[1, 2].G[2, 3].G[3, 1]
```

```
Out[•]= {{1, 0, 0}, {0, 0, 0}, {0, 0, 0}}
```

As claimed $G[1,1] = G[1, 2] \cdot G[2, 3] \cdot G[3, 1]$

For the 3-circuit starting with vertex 3, (3, 1, 2), the matrix has its one non zero term in row 3, column 3. The reader can show matrix for this product is $G[3,3]$.

```
In[•]:= MatrixForm[G[3, 1].G[1, 2].G[2, 3]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

For the 3-circuit starting with vertex 2, (2, 3, 1), the matrix has one non zero term in row 2, column 2.

```
In[•]:= MatrixForm[G[2, 3].G[3, 1].G[1, 2]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

These three NA matrices for the same simple 3 - circuit differ. Each has its own matrix. An incorrect order of these matrices gives 0 as the next example shows.

```
In[•]:= MatrixForm[G[2, 3].G[1, 2].G[3, 1]]
```

```
Out[•]//MatrixForm=
```

$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

Let $v[i,j]$ denote the value of the one-way arrow $a[i,j]$. The term $v[1,2]$ denotes the value of the arrow $a[1,2]$. The matrix showing the value of the round trip starting from vertex 1 and returning to vertex 1 is the following matrix $V[1]$. It has only one non term (1,1) on the diagonal.

$$V[1] := \begin{pmatrix} v[1, 2] * v[2, 3] * v[3, 1] & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

The second matrix is $V[2]$.

$$V[2] := \begin{pmatrix} 0 & 0 & 0 \\ 0 & v[2, 3] * v[3, 1] * v[1, 2] & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

The third matrix is $V[3]$

$$V[3] := \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & v[3, 1] * v[1, 2] * v[2, 3] \end{pmatrix}$$

Each $V[i]$ is a different matrix for the value of the different round trips of the simple 3-circuit. Although the round trips start from different vertexes, they have equal value shown by their matrices.

2. Super Circuits

Warning: Mathematica arranges the terms in the product in lexicographical order. This is incorrect. The sequence of arrows in the circuit is the correct order.

A super circuit is a set of simple circuits whose *arrows* do not intersect and share at least one *vertex*. The next two simple circuits

$\{(5,1),(1,4),(4,2),(2,5)\}$ and $\{(5,3),(3,4),(4,5)\}$

in the pentagon share no arrows and share vertex 5. They make the following super circuit. Its 5X5 matrix is $G[5,5]$

$\{(5,1),(1,4),(4,2),(2,5),(5,3),(3,4),(4,5)\}$.

The procedure step by step follows. To save space I replace G by g in the matrix.

```
In[ ]:= g14 := {{0, 0, 0, a[1, 4], 0}, {0, 0, 0, 0, 0},
              {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
In[ ]:= g53 := {{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0},
              {0, 0, 0, 0, 0}, {0, 0, a[5, 3], 0, 0}}
```

```
In[ ]:= g25 := {{0, 0, 0, 0, 0}, {0, 0, 0, 0, a[2, 5]},
              {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
In[ ]:= g31 := {{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0},
              {a[3, 1], 0, 0, 0, 0}, {0, 0, 0, 0, 0},
              {0, 0, 0, 0, 0}}
```

```
In[ ]:= g12 := {{a[1, 2], 0, 0, 0, 0}, {0, 0, 0, 0, 0},
              {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
In[*]:= g23 := {{0, 0, 0, 0, 0}, {0, 0, a[2, 3], 0, 0},
               {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
In[*]:= g34 := {{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0},
               {0, 0, 0, a[3, 4], 0}, {0, 0, 0, 0, 0},
               {0, 0, 0, 0, 0}}
```

```
In[*]:= g45 := {{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0},
               {0, 0, 0, 0, a[4, 5]}, {0, 0, 0, 0, 0}}
```

```
In[*]:= g51 := {{a[5, 1], 0, 0, 0, 0}, {0, 0, 0, 0, 0},
               {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
In[*]:= g42 := {{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0},
               {0, a[4, 2], 0, 0, 0}, {0, 0, 0, 0, 0}}
```

The next 3 steps make the following super circuit.

```
(g51.(g14.g42)).g25) (g53.g34).g45) =
((g51.(g14.g42)).g25).((g53.g34).g45)
```

Note the positions of the parentheses.

```
In[*]:= ((g51.(g14.g42)).g25)
```

```
{{0, 0, 0, 0, a[1, 4] * a[2, 5] * a[4, 2] * a[5, 1]},
 {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0},
 {0, 0, 0, 0, 0}}
```

```
In[*]:= (g53.g34).g45
```

```
{{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0},
 {0, 0, 0, 0, 0}, {0, 0, 0, 0, a[3, 4] * a[4, 5] * a[5, 3]}}
```

```
In[•]:= ((g51.(g14.g42)).g25).((g53.g34).g45)
```

```
{{0, 0, 0, 0, a[1, 4] * a[2, 5] * a[3, 4] * a[4, 2] *  
a[4, 5] * a[5, 1] * a[5, 3]}, {0, 0, 0, 0, 0},  
{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

This gives the super circuit in the first term on the diagonal of the matrix product $((g51.(g14.g42)).g25).((g53.g34).g45)$. Mathematica shows the factors in lexicographical order, not the correct super circuit order. Next is the procedure for $g23.g31 = g21$

```
In[•]:= g23.g31
```

```
{{0, 0, 0, 0, 0}, {a[2, 3] * a[3, 1], 0, 0, 0, 0},  
{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
g23.g31 = g21
```

```
In[•]:= (g23.g31).g12
```

```
{{0, 0, 0, 0, 0}, {a[1, 2] * a[2, 3] * a[3, 1], 0, 0, 0, 0},  
{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
(g23.g31).g12 = g21.g12 = g22
```

```
In[•]:= (g51.g14).g45
```

```
{{0, 0, 0, 0, a[1, 4] * a[4, 5] * a[5, 1]}, {0, 0, 0, 0, 0},  
{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

The final step is $(g51.g14).g45 = g54.g45$. It is the fifth term of the diagonal matrix $g[5,5]$.

```
In[•]:= ((g51.g14).g45)
```

```
{{0, 0, 0, 0, a[1, 4] * a[4, 5] * a[5, 1]}, {0, 0, 0, 0, 0},  
{0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}, {0, 0, 0, 0, 0}}
```

```
((g51.g14).g45).((g23.g31).g12)
```

$\{\{0, 0, 0, 0, 0\},$
 $\{0, 0, 0, 0, a[1, 2] * a[1, 4] * a[2, 3] * a[3, 1] * a[4, 5] * a[5, 1]\}, \{0, 0, 0, 0, 0\}, \{0, 0, 0, 0, 0\},$
 $\{0, 0, 0, 0, 0\}\}$

3. Valuation of a Circuit

The results in the preceding two sections show that the value of a circuit is given by its matrix whose terms are arrow values.

Let $1 + \rho_{i,j}$ denote the value of arrow $a[i, j]$. The value of the 3 – circuit, $V[\theta]$, is

$$V[\theta] = (1 + \rho_{12}) * (1 + \rho_{23}) * (1 + \rho_{31}) =$$

$$= 1 + \rho_{12} + \rho_{23} + \rho_{31} + \rho_{12} \rho_{23} + \rho_{12} \rho_{31} + \rho_{23} \rho_{31} + \rho_{12} \rho_{23} \rho_{31}.$$

The arithmetic mean of the value of the arrows in a circuit exceeds the geometric mean of their values unless $\rho_{ij} = \rho$. Given the mean, the geometric mean varies inversely with the dispersion of its terms. The closer together are $\rho_{i,j}$, the closer the geometric mean to the arithmetic mean. Given the arithmetic mean, the more nearly equal are $(1 + \rho_{i,j})$ in the circuit, the bigger the circuit value.

For $\rho_i \geq 0$,

$$(1) \quad \sum_1^n (1 + \rho_i) < \prod_1^n (1 + \rho_i).$$

The subscript identifies the owner of the arrow. Let w_i denote the return to owner i , $i = 1, 2, \dots, n$, $n = m(m-1)/2$.

Let $V(\theta)$ in (2) denote the value of circuit θ .

$$(2) \quad V(\theta) = \prod_1^n (1 + \rho_i).$$

The core constraint is inequality (3)

$$(3) \quad \sum_{i \in \theta} w_i \geq V(\theta)$$

If $w_i = 1 + \rho_i$, then the smallest acceptable payoff to owner i cannot satisfy (3) owing to inequality (1). The core constraints are not feasible. Changing to logs gives feasibility. Let the owner of vertex i get $\log w_i$ instead of w_i . The log of the value replaces value. Instead of (3), there is

$$(4) \quad \sum_{i \in \theta} \log w_i \geq \log V(\theta) = \log \prod_{i \in \theta} (1 + \rho_i) = \sum_{i \in \theta} \log(1 + \rho_i).$$

$\log w_i$ can satisfy (4) with equality even if some $\rho_i = 0$. Such arrows are unproductive. Their value = 1 and $\log 1 = 0$.

4 Undirected Edges

The undirected edge from vertex i to vertex j is denoted by $i \leftrightarrow j$. It differs from the pair of oppositely oriented arrows $i \rightarrow j$ and $j \rightarrow i$. The Nöther matrix for the undirected edge connecting vertex i and vertex j is

```
In[•]:= a := {{0, 1}, {1, 0}}
```

```
In[•]:= MatrixForm[a]
```

Out[•]//MatrixForm=

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

The matrix for the arrow $1 \rightarrow 2$ is a . The matrix for the arrow $2 \rightarrow 1$ is b

```
In[•]:= b := {{1, 0}, {0, 1}}
```

```
In[•]:= MatrixForm[b]
```

Out[•]//MatrixForm=

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

```
In[•]:= a.b
```

```
Out[•]= {{0, 1}, {1, 0}}
```

Thus $a.b = a$ and $\{1 \rightarrow 2, 2 \rightarrow 1\} = 1 \rightarrow 2$ as required.

```
In[•]:= Det[a.a]
```

```
Out[•]= 1
```

```
In[•]:= Transpose[a]
```

```
Out[•]= {{0, 1}, {1, 0}}
```

The transpose of a is b .

```
In[•]:= tl := {1 ↔ 2, 2 ↔ 3, 3 ↔ 4, 4 ↔ 2, 2 ↔ 1}
```

```
In[•]:= Array[0 &, {4, 4}]
```

```
Out[•]= {{0, 0, 0, 0}, {0, 0, 0, 0}, {0, 0, 0, 0}, {0, 0, 0, 0}}
```

```
In[•]:= g12 := {{0, 1, 0, 0}, {0, 0, 0, 0}, {0, 0, 0, 0},  
              {0, 0, 0, 0}}
```

```
In[•]:= g23 := {{0, 0, 0, 0}, {0, 0, 1, 0}, {0, 0, 0, 0},  
              {0, 0, 0, 0}}
```

```
In[•]:= g34 := {{0, 0, 0, 0}, {0, 0, 0, 0}, {0, 0, 0, 1},  
              {0, 0, 0, 0}}
```

```
In[•]:= g42 := {{0, 0, 0, 0}, {0, 0, 0, 0}, {0, 0, 0, 0},  
              {0, 1, 0, 0}}
```

```
In[•]:= g21 := {{0, 0, 0, 0}, {1, 0, 0, 0}, {0, 0, 0, 0},  
              {0, 0, 0, 0}}
```

```
In[•]:= g12.g23
```

```
In[•]:= g13 := {{0, 0, 1, 0}, {0, 0, 0, 0}, {0, 0, 0, 0},  
              {0, 0, 0, 0}}
```

```
In[•]:= g13.g34
```

```
In[•]:= g14 := {{0, 0, 0, 1}, {0, 0, 0, 0}, {0, 0, 0, 0},  
              {0, 0, 0, 0}}
```

```
In[•]:= g14.g42
```

```
In[•]:= g12 := {{0, 1, 0, 0}, {0, 0, 0, 0}, {0, 0, 0, 0},  
              {0, 0, 0, 0}}
```

```
In[•]:= g12.g21
```

```
In[•]:= g11 := {{1, 0, 0, 0}, {0, 0, 0, 0}, {0, 0, 0, 0},
               {0, 0, 0, 0}}
```

The Nöther matrix of tl is the product of the 5 matrices; $g12 \cdot g23 \cdot g34 \cdot g42 \cdot g21 = g11$, the circuit tl

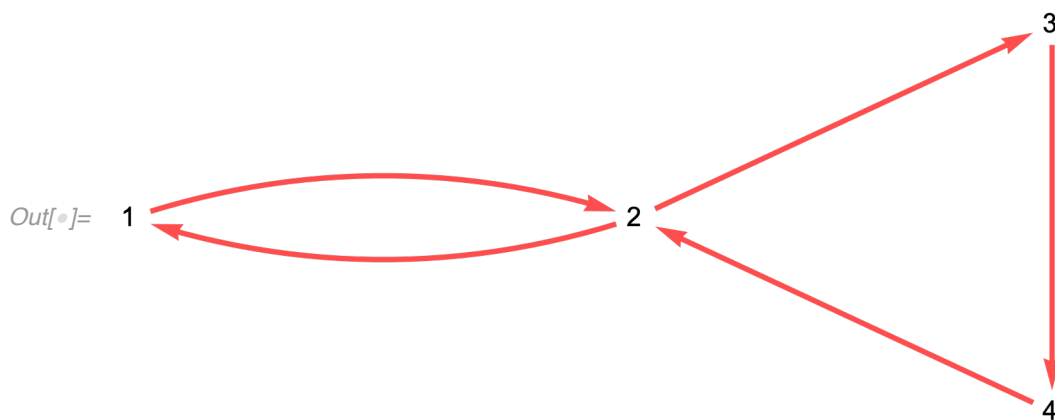
```
tl := {1 ↔ 2, 2 ↔ 3, 3 ↔ 4, 4 ↔ 2, 2 ↔ 1}
```

The adjacency matrix for tl is atl

```
In[•]:= atl := {{0, 1, 0, 0}, {1, 0, 1, 0}, {0, 0, 0, 1},
               {0, 1, 0, 0}}
```

```
In[•]:= gat1 := makeGrph[atl, 3]
```

```
In[•]:= gat1
```



The Undirected Edge $1 \leftrightarrow 2$ becomes the pair of oppositely Directed Edges $1 \rightarrow 2$ & $2 \rightarrow 1$ in the graph.

```
In[•]:= EulerianGraphQ[gat1]
```

```
Out[•]= True
```

```
In[•]:= HamiltonianGraphQ[gat1]
```

```
Out[•]= False
```

Therefore, an Eulerian graph can have two oppositely oriented arrows (Undirected Edge) made by two Directed Edges. This is impossible for a Hamiltonian graph.

Program